

**SE 2AA4, CS 2ME3 (Introduction to Software
Development)**

Winter 2018

23 Finite State Machines (Ch. 5)

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23 Finite State Machines (Ch. 5)

- Administrative details
- Classification of specification styles
- Continuation on specification qualities
- Homework exercise
- How to verify a specification
- Finite state machines

Administrative Details

- Some of today's slides adapted from Dr. Wassying's slides (and Ghezzi et al)
- A3
 - ▶ Part 1 - Specification: due 11:59 pm Mar 12
 - ▶ Part 2 - Code: due 11:59 pm Mar 26
 - ▶ spec.tex moved to se2aa4_cs2me3 repo
 - ▶ Minor corrections
 - ▶ LineT len output type better as \mathbb{N} (rather than \mathbb{Z})
 - ▶ LineT strt output should be *out* := PointT(s.x(), s.y())
 - ▶ PathT len output the total number of points (grid cells) on the path, including the beginning and end points (cells).
- A4
 - ▶ Your own design and specification
 - ▶ Due April 9 at 11:59 pm

A3

- make experiment
- make doc
- make test (uses catch2)
- Version 7 of g++
- Strongly recommend testing on mills
- On mills, modify .bashrc to include
 . /opt/rh/devtoolset-7/enable
- Use vectors for Seq2D constructor
- No pointers
- Follow tutorial approach for implementing templates

```
template <class T>
Seq2D<T>::Seq2D( vector<vector<T>> s, double scal
{ // details
}
```

Operational Versus Descriptive Spec

- Operational specification
 - ▶ “Let a be an array of n elements. The result of its sorting is an array b of n elements such that the first element of b is the minimum of a (if several elements of a have the same value, any one of them is acceptable); the second element of b is the minimum of the array of $n - 1$ elements obtained from a by removing its minimum element; and so on until all n elements of a have been removed.”
- Descriptive specification
 - ▶ What is the corresponding descriptive spec?

 - ▶ How can we further specify (formalize) the notion of sorted?

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 - ▶ What is the corresponding descriptive spec?
 - ▶ “The result of sorting array a is an array b which is a permutation of a and is sorted.”
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 - ▶ “Let a be an array of n elements. The result of its sorting is an array b of n elements such that the first element of b is the minimum of a (if several elements of a have the same value, any one of them is acceptable); the second element of b is the minimum of the array of $n - 1$ elements obtained from a by removing its minimum element; and so on until all n elements of a have been removed.”
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 - ▶ “The result of sorting array a is an array b which is a permutation of a and is sorted.”
 - ▶ How can we further specify (formalize) the notion of sorted?
 - ▶ $\text{sorted}(A) \equiv \forall(i : \mathbb{N} | 0 \leq i \leq (|A| - 2) : A[i] \leq A[i + 1])$

Homework Exercise

- Consider the **line formatter** specification and
 1. How well does the specification do with respect to the following qualities: abstract, correct, unambiguous, complete, consistent and verifiable?
 2. For a requirement specification like that given, what are the advantages and disadvantages of maintaining both a formal specification and a natural language specification?
- Even spending 5 minutes thinking about will help when we discuss next week
- In repo
 - ▶ The [line formatter specification](#)
 - ▶ [Meyer \(1985\)](#) “On Formalism in Specification”
- Will discuss next day

How to Verify a Specification

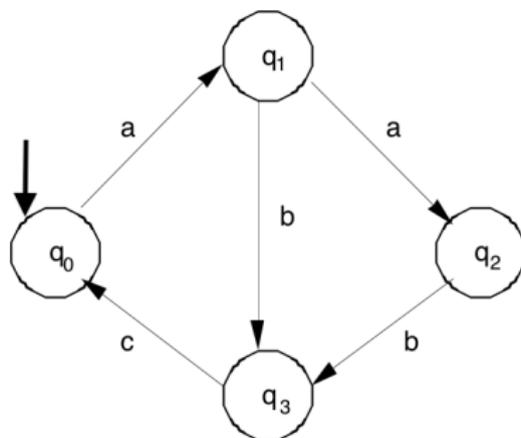
- How might you verify a specification?

How to Verify a Specification

- How might you verify a specification?
- *Observe* dynamic behaviour of the specified system
 - ▶ Simulation
 - ▶ Prototyping
 - ▶ “testing” the specification
- Mathematically analyze properties of the specified system, including proof
- Analogy with traditional engineering
 - ▶ Physical model of a bridge (prototype)
 - ▶ Mathematical model of a bridge
- We will return to this topic when we cover verification (Chapter 6)

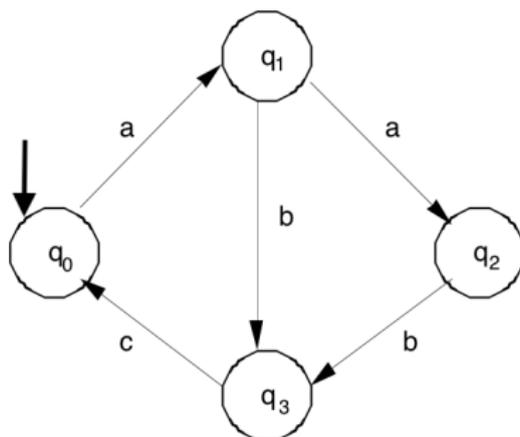
Finite State Machines (FSMs)

- Can specify control flow aspects
- Defined as
 - ▶ A finite set of states Q
 - ▶ A finite set of inputs I
 - ▶ A transition function $\delta : Q \times I \rightarrow Q$ (δ can be a partial function)



FSMs Continued

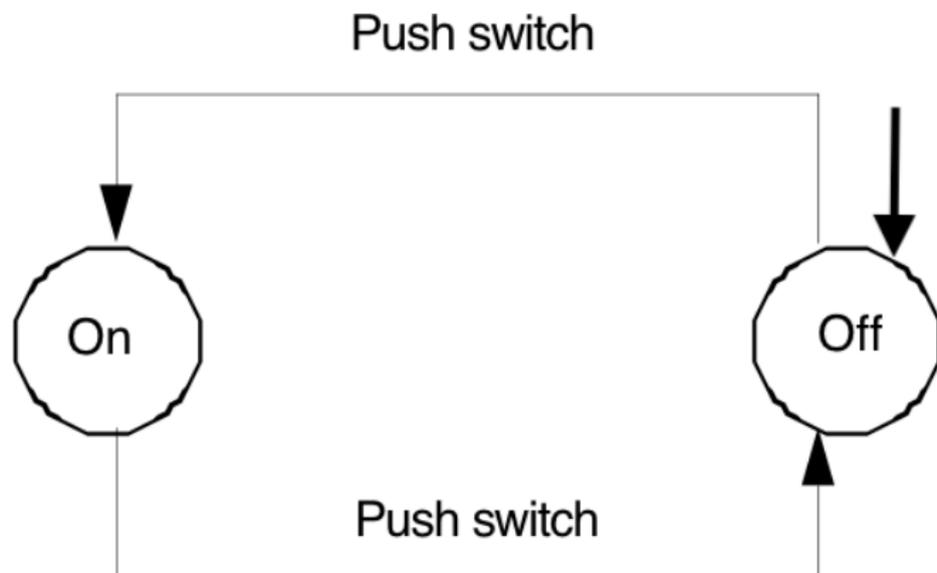
	q₀	q₁	q₂	q₃
a	<i>q₁</i>	<i>q₂</i>	-	-
b	-	<i>q₃</i>	<i>q₃</i>	-
c	-	-	-	<i>q₀</i>



Example: A Lamp

- What are the states Q for a typical lamp?
- What are the set of inputs I
- What is the transition function δ ?

Example: A Lamp

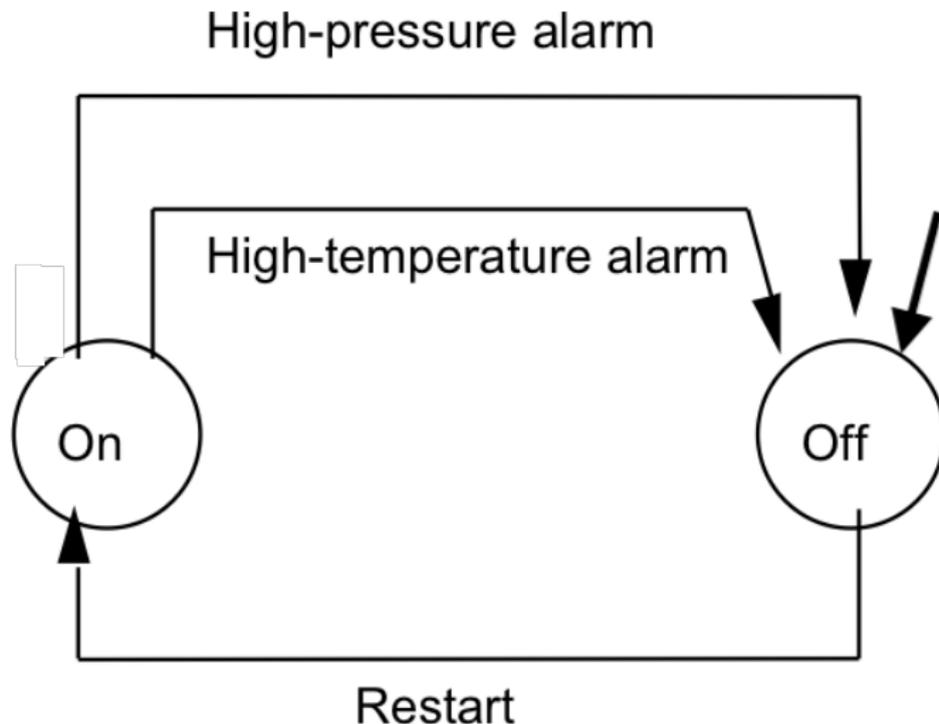


Example: A Plant Control System

Your plant is either On or Off. If an alarm occurs, what should be the state transition?

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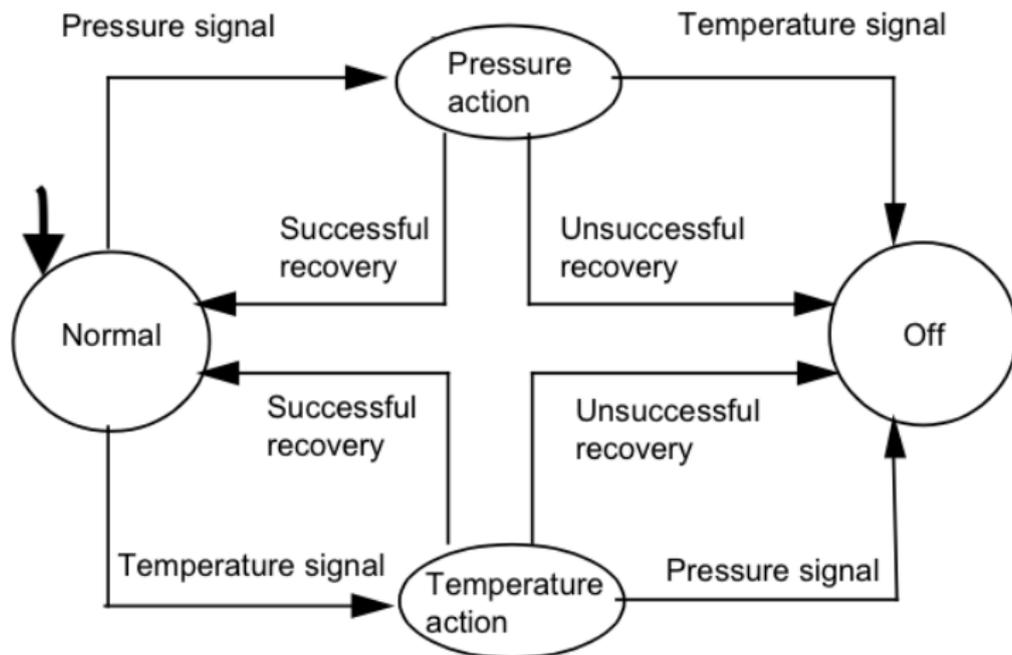


A Refinement

How might you refine the FSM if you have states for “pressure action” and “temperature action”?

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When to use FSMs for Specification?

- When is an FSM a good choice for specification?
- What are some examples of things we would specify using an FSM?

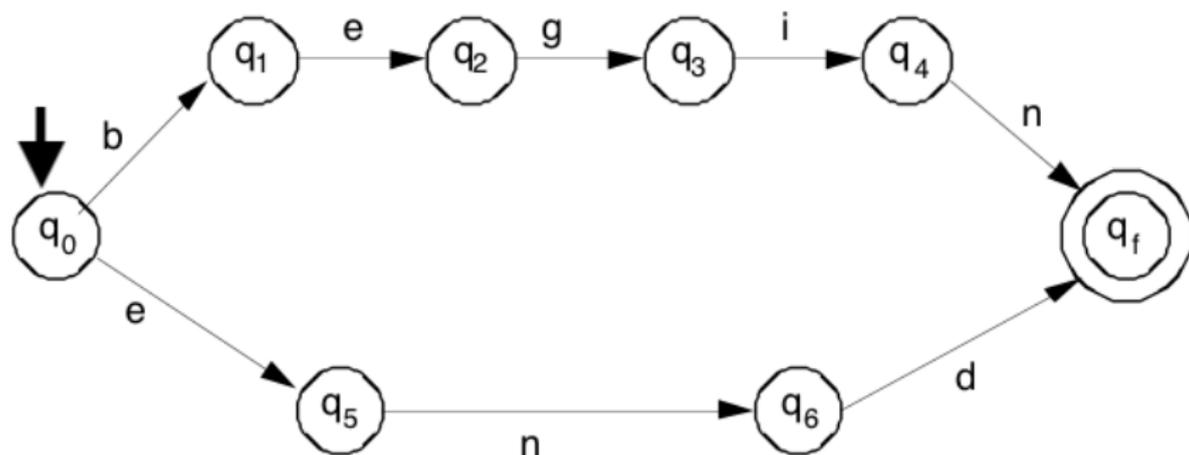
When to Potentially use FSMs

- Describing control flow
- Clear finite set of states (or modes)
- Specify acceptable strings for a parser
- Specifying hardware design
- For synchronous models (at any time a global state must be defined and a single transition must occur)

Classes of FSMs

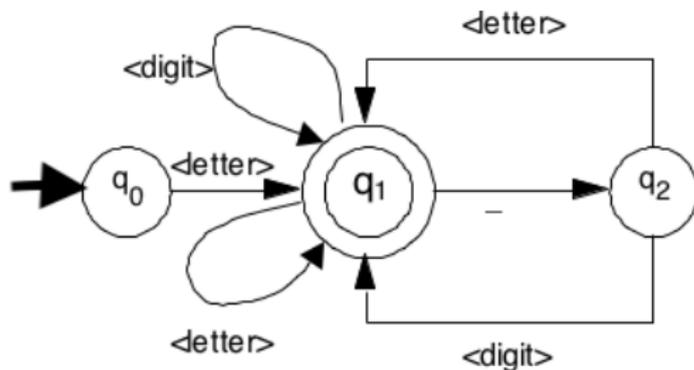
- Deterministic/nondeterministic
- FSMs as recognizers - introduce final states
- FSMs as transducers - introduce set of output
- ...

FSMs as Recognizers



What if an invalid character is entered?

FSMs as Recognizers Continued



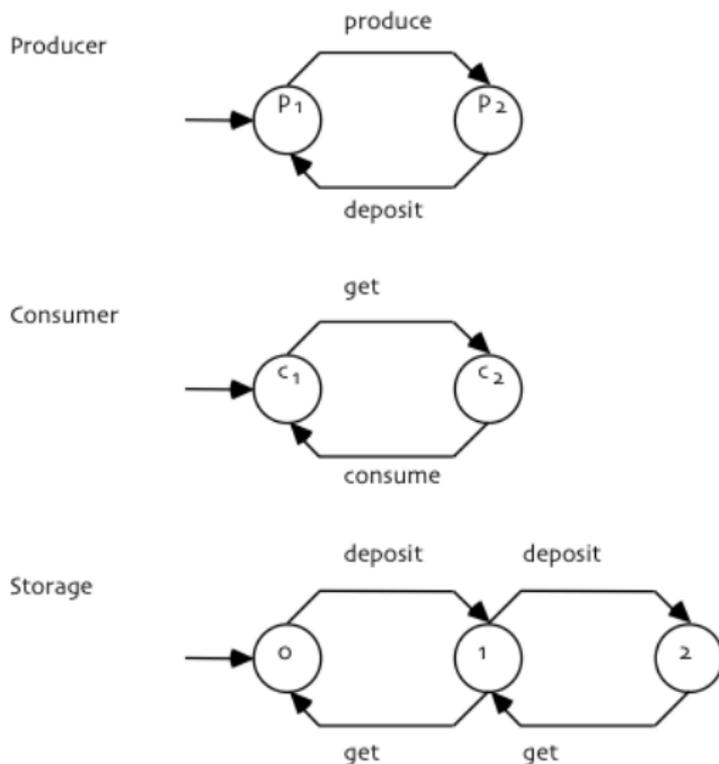
Legend: $\xrightarrow{\langle \text{letter} \rangle}$ is an abbreviation for a set of arrows
respectively

$\xrightarrow{\langle \text{digit} \rangle}$ is an abbreviation for a set of arrows
labeled 0, 1, ..., 9, respectively

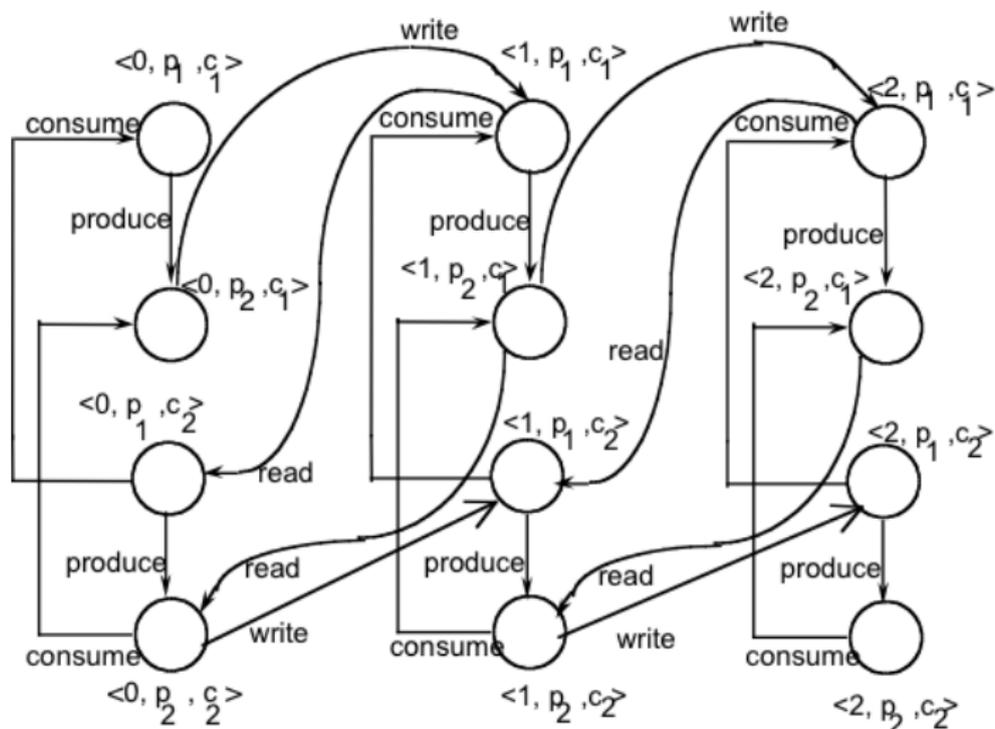
Limitations

- Finite memory
- State explosion - Given a number of FSMs with k_1, k_2, \dots, k_m states, their composition is an FSM with $k_1 \times k_2 \times \dots \times k_n$. This growth is exponential with the number of FSMs, not linear (we would like it to be $k_1 + k_2 + \dots + k_n$)

State Explosion: An Example



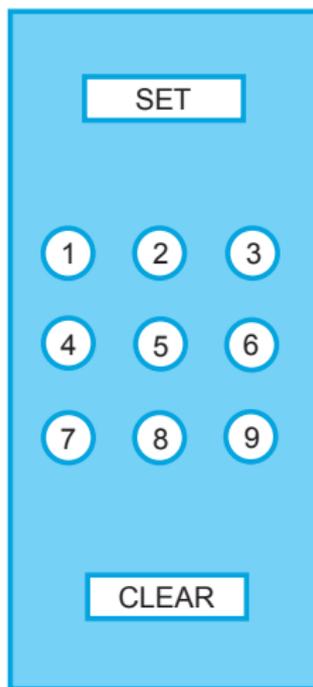
The Resulting FSM



Events Versus Conditions

- Events can be viewed as “pulses” in time - they do not last (retain their values)
- Conditions may retain their values indefinitely

FSM Example: Security Alarm



Security Alarm Example Continued

