

**SE 2AA4, CS 2ME3 (Introduction to Software  
Development)**

**Winter 2018**

## **11 Generic MIS (Ghezzi Ch. 4)**

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January 28, 2018



# 11 Generic MIS (Ghezzi Ch. 4)

- Administrative details
- ssh with X11 forwarding
- Homework exercise on specification
- Uses for abstract objects
- Implementing objects with other objects as state variables
- Generic modules
- Generic stack abstract object
- Properties exhibited by a stack module
- Generic queue ADT
- Access routine idioms
  - ▶ Set idioms
  - ▶ Sequence idioms
  - ▶ Tuple idioms
- Exceptions
- Quality Criteria

# Administrative Details

- Assignment 1
  - ▶ Partner files now in your repo
  - ▶ Part 2: January 31, 2018
  - ▶ If you know how, tag your assignment submissions
- Questions?

# X Windows

- If you want to use a gui on mills, you can use X11 forwarding
- On a Mac or with Linux, you use ssh -X or ssh -Y
- ssh -X mills.cas.mcmaster.ca (or -Y)
- On a Mac you need XQuartz installed
- If you know the details for Windows, please post to Avenue

# Homework Answer: Access Routine Semantics

totalArea():

- output

$$out := ? \quad (i : \mathbb{N} | i \in [0..|s| - 1] : s[i].area())$$

averageRadius():

- output

$$out :=$$

# Homework Answer: Access Routine Semantics

totalArea():

- output

$$out := + (i : \mathbb{N} | i \in [0..|s| - 1] : s[i].area())$$

averageRadius():

- output

$$out :=$$

# Homework Answer: Access Routine Semantics

totalArea():

- output

$$out := + (i : \mathbb{N} | i \in [0..|s| - 1] : s[i].area())$$

or?

averageRadius():

- output

$$out :=$$

# Homework Answer: Access Routine Semantics

totalArea():

- output

$$out := + (i : \mathbb{N} | i \in [0..|s|-1] : s[i].area())$$

or?

$$out := + (c : CircleT | c \in s : c.area())$$

averageRadius():

- output

$$out := ?$$

# Homework Answer: Access Routine Semantics

totalArea():

- output

$$out := + (i : \mathbb{N} | i \in [0..|s| - 1] : s[i].area())$$

or?

$$out := + (c : CircleT | c \in s : c.area())$$

averageRadius():

- output

$$out := + (c : CircleT | c \in s : c.radius()) / |s|$$

# Uses for Abstract Objects

- Creating a single abstract object corresponds to the singleton design pattern
- Provides “global” variables
- Uses
  - ▶ Shared resource
  - ▶ Hoffman and Strooper example (assembler)
  - ▶ Look up table for global state (read only)
  - ▶ Logger (write only)
- Problematic for testing if high coupling and frequent state changes, since many test cases will depend on the state of the abstract object

# Objects with Other Objects as State Variables

Potential prob with the naive implement of deque of CircleT?

```
@staticmethod  
def pushBack(c):  
    Deq.s = Deq.s + [c]
```

Consider

```
p1 = PointT(1.5, 2)  
c1 = CircleT(p1, 5)  
Deq.init()  
Deq.pushBack(c1)  
c1.r = 6 #circumventing information hiding  
print(c1.rad())  
print(Deq.back().rad())
```

# Solutions to Potential Problem

- Interface prevents potential abuse
  - ▶ Provide no mutators in the interface, just a constructor and selectors
  - ▶ Assume information hiding will be respected
- A more robust implementation
  - ▶ The state variable stores copies of objects, not references
  - ▶ Use copy library in Python

```
import copy  
...  
@staticmethod  
def pushBack(c):  
    cnew = copy.deepcopy(c)  
    Deq.s = Deq.s + [cnew]
```

## New Problem for Robust Implementation

```
p1 = PointT(1.5, 2)
c1 = CircleT(p1, 5)
Deq.init()
Deq.pushBack(c1)
print(c1 == Deq.back())
```

Why is the behaviour not what we would naively expect?

What can we do about it?

## New Problem for Robust Implementation

```
p1 = PointT(1.5, 2)
c1 = CircleT(p1, 5)
Deq.init()
Deq.pushBack(c1)
print(c1 == Deq.back())
```

Why is the behaviour not what we would naively expect?

What can we do about it?

## Redefine Equality for CircleT

```
class CircleT:  
  
    def __init__(self, cin, rin):  
        self.c = cin  
        self.r = rin  
  
    def __eq__(self, cin):  
        return self.c == cin.cen() and  
            self.r == cin.rad()
```

What other classes will need \_\_eq\_\_ redefined?

## Examples of A2 from 2017

- Solution assuming information hiding in repo of previous assignment solutions ([here](#))
- Robust solution in the same folder as these lecture slides ([here](#))

# Generic Modules

- What if we have a sequence of integers, instead of a sequence of point masses?
- What if we want a stack of integers, or characters, or pointT, or pointMassT?
- Do we need a new specification for each new abstract object?
- No, we can have a single abstract specification implementing a family of abstract objects that are distinguished only by a few variabilities
- Rather than duplicate nearly identical modules, we parameterize one **generic module** with respect to type(s)
- Advantages
  - ▶ Eliminate chance of inconsistencies between modules
  - ▶ Localize effects of possible modifications
  - ▶ Reuse

# Generic Stack Module Syntax

## Generic Module

Stack(T)

## Exported Constants

MAX\_SIZE = 100

## Exported Access Programs

Routine name	In	Out	Exceptions
...	...	...	...

# Stack Module Syntax

## Exported Access Programs

Routine name	In	Out	Exceptions
s_init			
s_push	T		FULL
s_pop			EMPTY
s_top		T	EMPTY
s_depth		integer	

# Semantics

## **State Variables**

## **State Invariant**

## **Assumptions**

# Semantics

## **State Variables**

$s$ : sequence of  $T$

## **State Invariant**

## **Assumptions**

# Semantics

## State Variables

$s$ : sequence of  $T$

## State Invariant

$|s| \leq \text{MAX\_SIZE}$

## Assumptions

# Semantics

## State Variables

$s$ : sequence of  $T$

## State Invariant

$|s| \leq \text{MAX\_SIZE}$

## Assumptions

`s_init()` is called before any other access routine

# Access Routine Semantics

s\_init():

- transition:
- exception:

s\_push(x):

- transition:
- exception:

s\_pop():

- transition:
- exception:

# Access Routine Semantics

s\_init():

- transition:  $s := <>$
- exception:

s\_push(x):

- transition:
- exception:

s\_pop():

- transition:
- exception:

# Access Routine Semantics

s\_init():

- transition:  $s := <>$
- exception: none

s\_push(x):

- transition:
- exception:

s\_pop():

- transition:
- exception:

# Access Routine Semantics

s\_init():

- transition:  $s := <>$
- exception: none

s\_push(x):

- transition:  $s := s || < x >$
- exception:

s\_pop():

- transition:
- exception:

# Access Routine Semantics

s\_init():

- transition:  $s := <>$
- exception: none

s\_push(x):

- transition:  $s := s || < x >$
- exception:  $exc := (|s| = \text{MAX\_SIZE} \Rightarrow \text{FULL})$

s\_pop():

- transition:
- exception:

# Access Routine Semantics

s\_init():

- transition:  $s := <>$
- exception: none

s\_push(x):

- transition:  $s := s || < x >$
- exception:  $exc := (|s| = \text{MAX\_SIZE} \Rightarrow \text{FULL})$

s\_pop():

- transition:  $s := s[0..|s| - 2]$
- exception:

# Access Routine Semantics

s\_init():

- transition:  $s := <>$
- exception: none

s\_push(x):

- transition:  $s := s || < x >$
- exception:  $exc := (|s| = \text{MAX\_SIZE} \Rightarrow \text{FULL})$

s\_pop():

- transition:  $s := s[0..|s| - 2]$
- exception:  $exc := (|s| = 0 \Rightarrow \text{EMPTY})$

# Access Routine Semantics Continued

s\_top():

- output:
- exception:

s\_depth():

- output:
- exception:

# Access Routine Semantics Continued

s\_top():

- output:  $out := s[|s| - 1]$
- exception:

s\_depth():

- output:
- exception:

# Access Routine Semantics Continued

s\_top():

- output:  $out := s[|s| - 1]$
- exception:  $exc := (|s| = 0 \Rightarrow \text{EMPTY})$

s\_depth():

- output:
- exception:

# Access Routine Semantics Continued

s\_top():

- output:  $out := s[|s| - 1]$
- exception:  $exc := (|s| = 0 \Rightarrow \text{EMPTY})$

s\_depth():

- output:  $out := |s|$
- exception:

# Access Routine Semantics Continued

s\_top():

- output:  $out := s[|s| - 1]$
- exception:  $exc := (|s| = 0 \Rightarrow \text{EMPTY})$

s\_depth():

- output:  $out := |s|$
- exception: **none**

# Stack Module Properties

$\{true\}$

s\_init()

$\{|s'| = 0\}$

$\{|s| < \text{MAX\_SIZE}\}$

s\_push( $x$ )

$\{|s'| = |s| + 1 \wedge s' [|s'| - 1] = x \wedge s'[0..|s| - 1] = s[0..|s| - 1]\}$

$\{|s| < \text{MAX\_SIZE}\}$

s\_push( $x$ )

s\_pop()

$\{s' = s\}$

# Generic Queue Module ADT Syntax

## Generic Template Module

QueueADT(T)

### Exported Types

QueueT = ?

### Exported Constants

MAX\_SIZE = 100

### Exported Access Programs

Routine name	In	Out	Exceptions
...	...	...	...

# QueueADT Module Syntax

## Exported Access Programs

Routine name	In	Out	Exceptions
new QueueT		QueueT	
add	T		queue_full
pop			queue_empty
front		T	queue_empty
isempty		boolean	

# Semantics

## **State Variables**

## **State Invariant**

## **Assumptions**

# Semantics

## **State Variables**

$s$ : sequence of  $T$

## **State Invariant**

## **Assumptions**

# Semantics

## State Variables

$s$ : sequence of  $T$

## State Invariant

$|s| \leq \text{MAX\_SIZE}$

## Assumptions

# Access Routine Semantics

new QueueT():

- transition:
- output:
- exception:

add(x):

- transition:
- exception:

pop():

- transition:
- exception:

# Access Routine Semantics

new QueueT():

- transition:  $s := <>$
- output:  $out := self$
- exception: none

add(x):

- transition:
- exception:

pop():

- transition:
- exception:

# Access Routine Semantics

new QueueT():

- transition:  $s := <>$
- output:  $out := self$
- exception: none

add(x):

- transition:  $s := s || < x >$
- exception:  $exc := (|s| = \text{MAX\_SIZE} \Rightarrow \text{queue\_full})$

pop():

- transition:
- exception:

# Access Routine Semantics

new QueueT():

- transition:  $s := <>$
- output:  $out := self$
- exception: none

add(x):

- transition:  $s := s || < x >$
- exception:  $exc := (|s| = \text{MAX\_SIZE} \Rightarrow \text{queue\_full})$

pop():

- transition:  $s := s[1..|s| - 1]$
- exception:  $exc := (|s| = 0 \Rightarrow \text{queue\_empty})$

# Access Routine Semantics Continued

front():

- output:
- exception:

isempty():

- output:
- exception:

# Access Routine Semantics Continued

front():

- output:  $out := s[0]$
- exception:  $exc := (|s| = 0 \Rightarrow \text{queue\_empty})$

isempty():

- output:
- exception:

# Access Routine Semantics Continued

front():

- output:  $out := s[0]$
- exception:  $exc := (|s| = 0 \Rightarrow \text{queue\_empty})$

isempty():

- output:  $out := |s| = 0$
- exception: **none**

# Queue Module Properties

$\{true\}$

q.init()

$\{|s'| = 0 \wedge is\_init\}$

$\{|s| < MAX\_SIZE\}$

add( $x$ )

$\{|s'| = |s| + 1 \wedge s'[0] = x \wedge s'[1..|s'| - 1] = s[0..|s| - 1]\}$

## Set Idiom

Routine name	In	Out	Exceptions
set_add	T		Member, Full
set_del	T		NotMember
set_member	T	boolean	
set_size		integer	

# Sequence Idiom

Routine name	In	Out	Exceptions
seq_init			
seq_add	integer, T		PosOutOfRange, Full
seq_del	integer		PosOutOfRange
seq_setval	integer, T		PosOutOfRange
seq_getval	integer	T	PosOutOfRange
seq_size		integer	
seq_start			
seq_next		T	AtEnd
seq_end		boolean	
seq_append	T		Full

# Tuple Idiom Version 1

Routine name	In	Out	Exceptions
tp_init			
tp_set_f <sub>1</sub>	T <sub>1</sub>		
tp_get_f <sub>1</sub>		T <sub>1</sub>	
...	...	...	...
tp_set_f <sub>N</sub>	T <sub>N</sub>		
tp_get_f <sub>N</sub>		T <sub>N</sub>	

## Tuple Idiom Version 2

Routine name	In	Out	Exceptions
tp_init			
tp_set	$T_1, T_2, \dots, T_N$		
tp_get		T	

## Example Subclass Exception in Python

```
class Full(Exception):
    def __init__(self, value):
        self.value = value
    def __str__(self):
        return str(self.value)
```

Example of raising the exception

```
if size == Seq.MAX_SIZE:
    raise Full("Maximum size exceeded")
```

# Exception Signaling

- Useful to think about exceptions in the design process
- Will need to decide how exception signalling will be done
  - ▶ A special return value, a special status parameter, a global variable
  - ▶ Invoking an exception procedure
  - ▶ Using built-in language constructs
- Caused by errors made by programmers, not by users
- Write code so that it avoid exceptions
- Exceptions will be particularly useful during testing

# Assumptions versus Exceptions

- The assumptions section lists assumptions the module developer is permitted to make about the programmer's behaviour
- Assumptions are expressed in prose
- Use assumptions to simplify the MIS and to reduce the complexity of the final implementation
- Interface design should provide the programmer with a means to check so that they can avoid exceptions
- When an exception occurs no state transitions should take place, any output is *don't care*

# Quality Criteria

- Consistent
  - ▶ Name conventions
  - ▶ Ordering of parameters in argument lists
  - ▶ Exception handling, etc.
- Essential - omit unnecessary features
- General - cannot always predict how the module will be used
- As implementation independent as possible
- Minimal - avoid access routines with two potentially independent services
- High cohesion - components are closely related
- Low coupling - not strongly dependent on other modules
- Opaque - information hiding